

Almost Affine Lambda Terms

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Jaśkowski, S. 1963.
Über Tautologien, in welchen keine Variable
mehr als zweimal vorkommt.
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der Mathematik* **9** (12–15), 219–228.

ÜBER TAUTOLOGIEN, IN WELCHEN KEINE VARIABLE
MEHR ALS ZWEIMAL VORKOMMT
von S. JAŚKOWSKI in Toruń (Polen)

H. THIELE hat im Jahre 1960 das Problem gestellt, das implikative Teilsystem
des Aussagenkalküls mit den Axiomen

$CCpqCCqCp$, $CCpCqCq$, $CpCqp$

zu untersuchen. Hier wird für dieses System¹⁾ und für ein anderes, in dem das
letzte Axiom durch ein schwächeres, nämlich durch Cqp ersetzt wird, ein Entschei-
dungsverfahren angegeben. Die Methode beruht auf einer Untersuchung von ge-
wissen allgemeinen Eigenschaften der Ausdrücke, in welchen keine Satzvariable
mehr als zweimal vorkommt. Dabei wird eine dreiwertige Matrix benutzt. Das
Problem der charakteristischen Matrizen für die oben erwähnten Systeme ist bisher
nicht gelöst.

1. Ein Lemma über Äquivalenzrelationen

Reflexive, symmetrische und transitive Relationen auf einer Menge M heißen
Äquivalenzrelationen über M . Es seien \sim_1, \sim_2 Äquivalenzrelationen über der
Menge M . Es ist bekannt, daß es über M eine eindeutig bestimmte Kompositions-
äquivalenz \sim_3 von \sim_1, \sim_2 mit folgenden Eigenschaften gibt:

On tautologies, in which no
[propositional] variable occurs more
than twice.

Hirokawa, S. 1992.
Balanced formulas, BCK-minimal formulas and
their proofs.
Logical Foundations of Computer Science — Tver '92,
pp. 198–208.

Here a formula is balanced iff no type variable occurs more than twice. Balanced
formulas are called as formulas with one-two-property in [6, 12]. We denote the set
of balanced formulas as $F_{1,2}$. The author learned a result by Jaskowski [12] from P.
Idzjak which states that

$BCK \cap F_{1,2} = LJ \cap F_{1,2}$.

By analysis of the type assignment figures, we prove that if a λ -term in β -normal form
has balanced type, then it is a BCK- λ -term. This gives a direct proof for Jaskowski's
result.

A balanced sequent provable in intuitionistic logic has
an *affine* inhabitant.

Hirokawa slightly improved
Jaskowski's result.

λ-terms

| | |
|---------------|-------------|
| x | variable |
| $\lambda x.M$ | abstraction |
| MN | application |

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Affine λ-terms

| | |
|---------------|--|
| x | variable |
| $\lambda x.M$ | abstraction |
| MN | application $FV(M) \cap FV(N) = \emptyset$ |

Affine λ-terms are *typable* (Hindley 1989).

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Every variable is used at most once.
Interesting properties of affine lambda terms have to do with their typings.

Type assignment to λ-terms

$x_1:\alpha_1, \dots, x_n:\alpha_n \Rightarrow M:\alpha$
type environment

$x:\alpha \Rightarrow x:\alpha$

$$\frac{\Gamma \Rightarrow M:\beta}{\Gamma - \{x:\alpha\} \Rightarrow \lambda x.M:\alpha \rightarrow \beta} \rightarrow I \quad \frac{\Gamma \Rightarrow M:\alpha \rightarrow \beta \quad \Delta \Rightarrow N:\alpha}{\Gamma \cup \Delta \Rightarrow MN:\beta} \rightarrow E$$

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Type environment includes type declarations for all and only variables that occur free in the term.

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Some terminology.

$$\Gamma \Rightarrow \alpha$$

sequent

- $\vdash \Gamma \Rightarrow M:\alpha$
- $\Gamma \Rightarrow \alpha$ is a **typing** of M
- M is an **inhabitant** of $\Gamma \Rightarrow \alpha$

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$$\Gamma \Rightarrow \alpha \text{ is a } \text{principal} \text{ typing of } M$$


$$\vdash \Gamma' \Rightarrow M:\alpha' \text{ just in case } (\Gamma' \Rightarrow \alpha') = (\Gamma \Rightarrow \alpha)\theta$$

for some type substitution θ

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Polarity of type occurrences

$$x:(\overset{-}{p} \rightarrow \overset{+}{q}) \rightarrow \overset{+}{r} \rightarrow s, y:\overset{-}{p} \rightarrow \overset{+}{q} \Rightarrow \overset{-}{r} \rightarrow \overset{+}{s}$$

$$\Gamma \Rightarrow \alpha \text{ is } \text{balanced}$$


each atomic type has at most 1 positive and
at most 1 negative occurrence in $\Gamma \Rightarrow \alpha$

This stricter definition of
“balanced” is from Mints’ textbook
on intuitionistic logic and is more
convenient for my purposes.

Affine λ -terms and balanced sequents

- Principal typing of an affine λ -term is balanced (Belnap 1976, Hirokawa 1992).
- A balanced sequent has at most one inhabitant up to $\beta\eta$ -equality (**Coherence Theorem**, Mints 1981, Babaev and Solov'ev 1982).
- A β -normal inhabitant of a balanced sequent is always affine (Jaśkowski 1963, Hirokawa 1992).

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These properties extend to what I call “almost affine” lambda terms.

Almost affine λ -terms

$$\begin{array}{c}
 x:\alpha \Rightarrow x:\alpha \\
 \frac{\Gamma \Rightarrow M:\beta}{\Gamma - \{x:\alpha\} \Rightarrow \lambda x.M:\alpha \rightarrow \beta} \rightarrow I \quad \frac{\Gamma \Rightarrow M:\alpha \rightarrow \beta \quad \Delta \Rightarrow N:\alpha}{\Gamma \cup \Delta \Rightarrow MN:\beta} \rightarrow E \\
 \text{ran}(\Gamma \cap \Delta) \subseteq \text{At}
 \end{array}$$

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The definition of almost affine lambda-terms refer to their typing.

$$\frac{\frac{\frac{y:q \rightarrow q \rightarrow p \quad x:q}{yx:q \rightarrow p} \rightarrow E \quad x:q}{yxx:p} \rightarrow E \quad \frac{z:r \rightarrow q \quad w:r}{zw:q} \rightarrow E}{\lambda x.yxx:q \rightarrow p} \rightarrow I \quad \frac{\lambda x.yxx:q \rightarrow p \quad zw:q}{(\lambda x.yxx)(zw):p} \rightarrow E$$

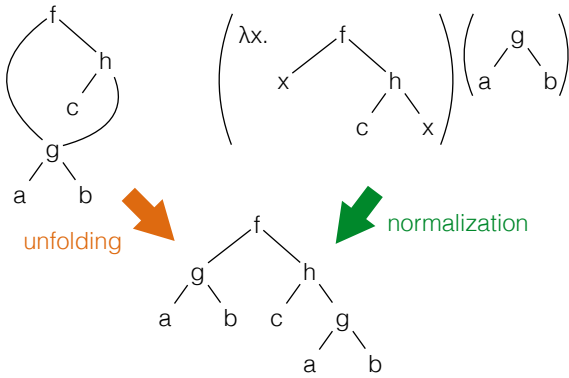
$$\frac{\frac{y:q \rightarrow q \rightarrow p \quad \frac{z:r \rightarrow q \quad w:r}{zw:q} \rightarrow E}{y(zw):q \rightarrow p} \rightarrow E \quad \frac{z:r \rightarrow q \quad w:r}{zw:q} \rightarrow E}{y(zw)(zw):p} \rightarrow E$$

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An example of an almost affine lambda term.
Its beta-normal form is not almost affine.

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Term graph (tree with sharing)



Almost affine lambda terms may be thought of as a generalization of “term graphs”.

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$$y:q \overset{+}{\rightarrow} q \overset{+}{\rightarrow} p, z:r \overset{-}{\rightarrow} q, w:r \Rightarrow (\lambda x. yxx)(zw):p \overset{+}{\rightarrow}$$

$\Gamma \Rightarrow \alpha$ is **negatively non-duplicated**



each atomic type has at most one **negative** occurrence in $\Gamma \Rightarrow \alpha$

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Almost affine λ -terms and negatively non-duplicated sequents

- Principal typing of an almost affine λ -term is negatively non-duplicated (Aoto 1999, Kanazawa 2007, 2011).
- A negatively non-duplicated sequent has at most one inhabitant up to $\beta\eta$ -equality (Aoto and Ono 1994).
- An inhabitant of a negatively non-duplicated sequent is always $\beta\eta$ -equal to an almost affine λ -term.

The coherence theorem about affine lambda terms and the correspondence between affine lambda terms and balanced sequents extend to almost affine lambda terms and negatively non-duplicated sequents.

This paper proves the third bullet point.

Affine λ -terms and balanced sequents

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Affine vs. balanced.

- Principal typing of an **affine** λ -term is **balanced**.
- A **balanced** sequent has at most one inhabitant up to $\beta\eta$ -equality.
- A β -normal inhabitant of a **balanced** sequent is always **affine**.

Almost affine λ -terms and negatively non-duplicated sequents

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Almost affine vs. negatively non-duplicated.

The third property is stated slightly differently due to the fact that the class of almost affine lambda-terms is not closed under beta-reduction.

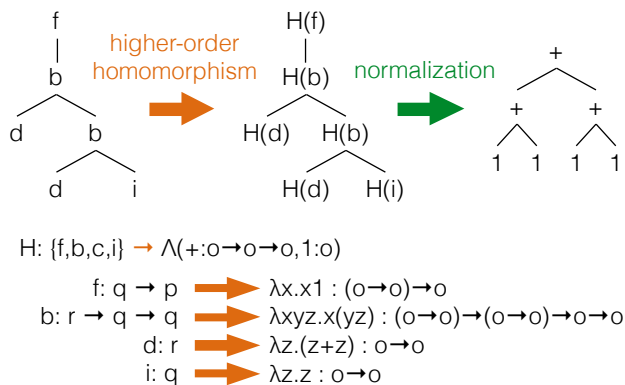
- Principal typing of an **almost affine** λ -term is **negatively non-duplicated**.
- A **negatively non-duplicated** sequent has at most one inhabitant up to $\beta\eta$ -equality.
- An inhabitant of a **negatively non-duplicated** sequent is always $\beta\eta$ -equal to an **almost affine** λ -term.

Almost affine λ -terms and tree transductions

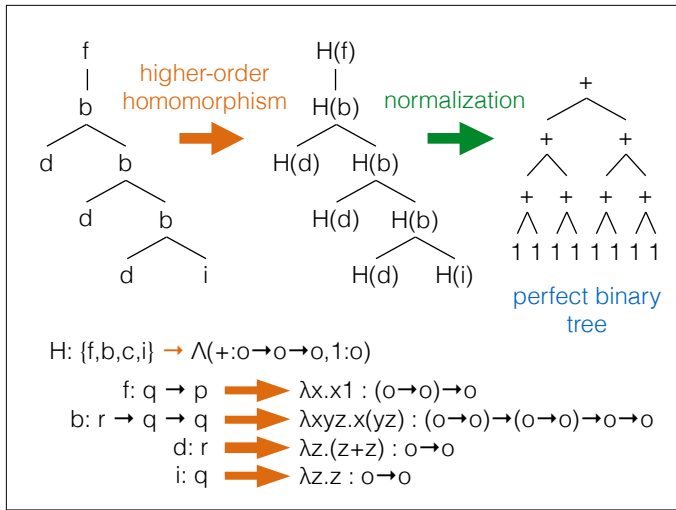
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Why would one be interested in the question of what lambda terms inhabit negatively non-duplicated sequents?

Almost affine lambda terms characterize an important class of tree transformations.

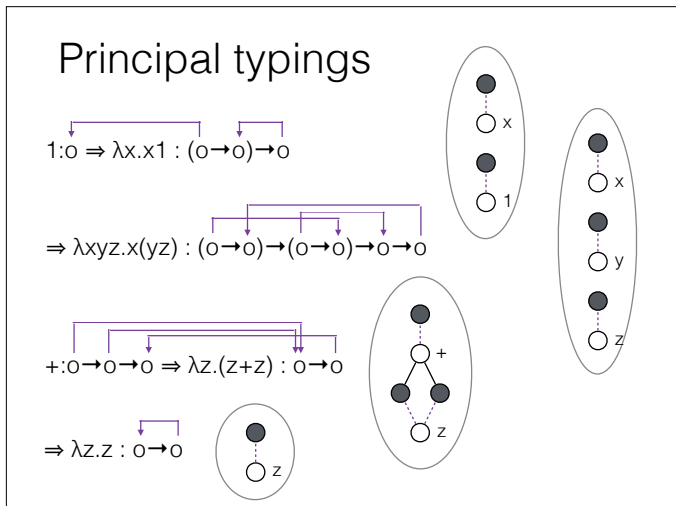


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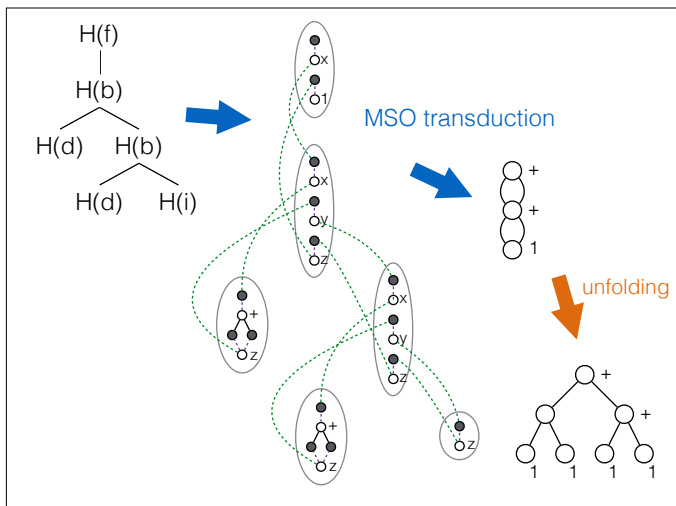
With almost affine lambda terms, normalization is easy to carry out (a kind of normalization by evaluation) using principal typings.

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Graph representations of principal typings.

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We can compute the normal form using these graphs.

The output graph is defined in terms of relations on the input graph defined by MSO formulas.

Tree transductions

Almost affine higher-order
homomorphism
+
normalization

\equiv

MSO transduction
+
unfolding

Kanazawa 2009

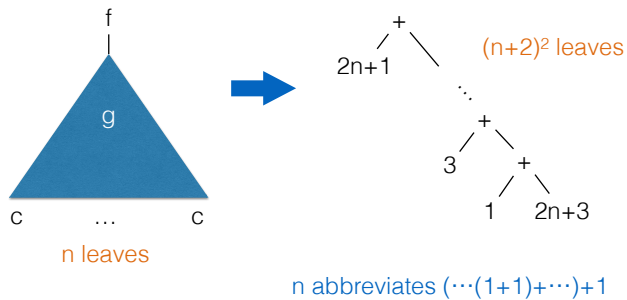
- Principal typing as graph
- Principal typing of an almost affine λ -term is negatively non-duplicated

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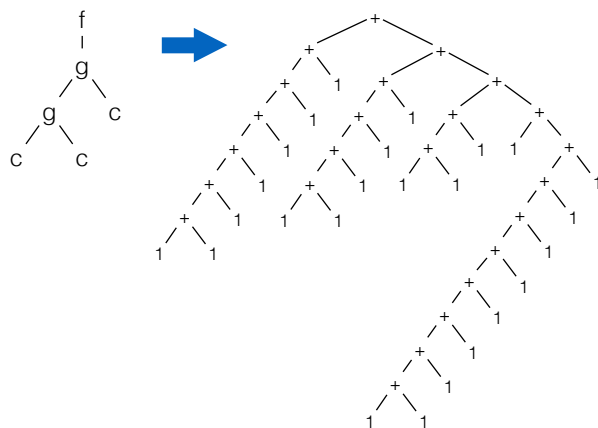
Correspondence.

One direction holds for all lambda terms that have negatively non-duplicated principal typings.

A more complex example

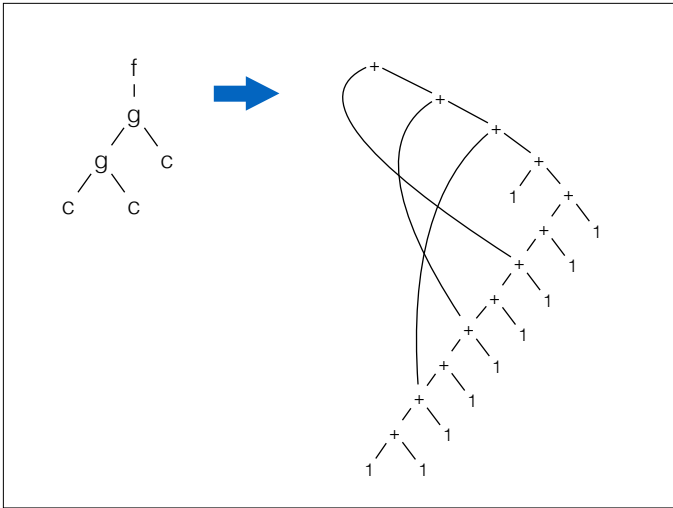


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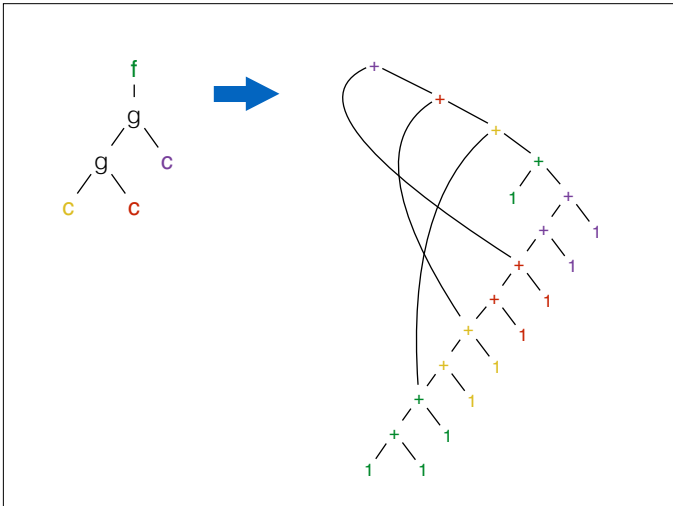


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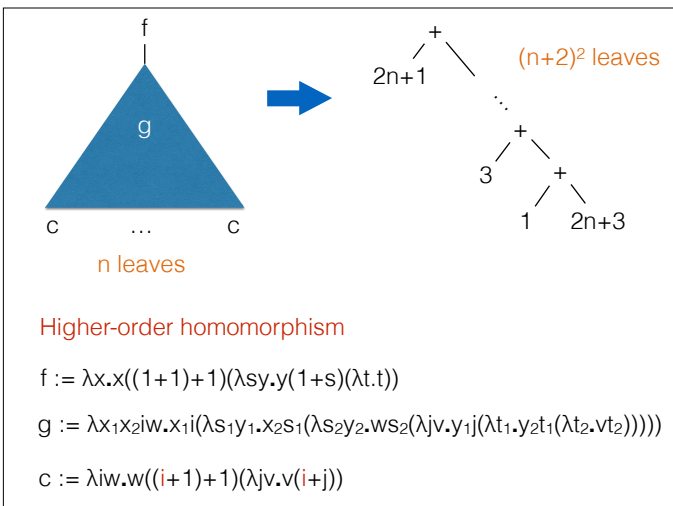


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This term graph can be obtained by MSO transduction.

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So there's a corresponding almost affine higher-order homomorphism.

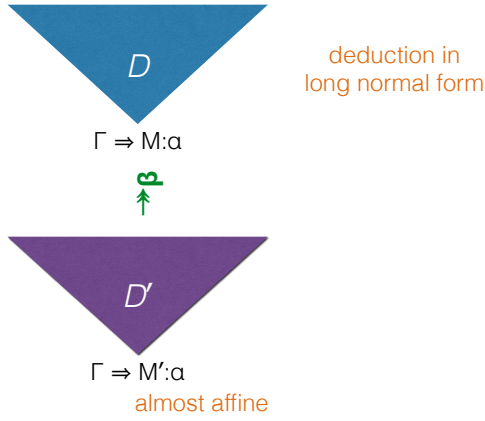
Almost affine λ -terms and negatively non-duplicated sequents

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MSO transduction is obtained only for almost affine higher-order homomorphism.

$\Gamma \Rightarrow \alpha$ negatively non-duplicated



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Some idea of how the theorem is proved.

$$\frac{y:\beta_1 \rightarrow \dots \rightarrow \beta_n \rightarrow p \quad \Gamma_1 \Rightarrow M_1:\beta_1 \quad \dots \quad \Gamma_n \Rightarrow M_n:\beta_n}{y:\beta_1 \rightarrow \dots \rightarrow \beta_n \rightarrow p, \Gamma_1 \cup \dots \cup \Gamma_n \Rightarrow yM_1 M_n:p} \rightarrow E$$

$$\{x:\delta \in \Gamma \mid \delta \notin \text{At}, x:\delta \in \Gamma_i \cap \Gamma_j, i \neq j\} = \{y_1:\delta_1, \dots, y_m:\delta_m\} \neq \emptyset$$

$$\text{tail}(\delta_i) = q_i$$

Find a k such that
 $M \equiv_{\alpha} N[z := y_k P_1 \dots P_l]$ and
 $x \in \text{FV}(N) \cap \text{FV}(y_k P_1 \dots P_l)$ implies $\Gamma(x) \in \text{At}$

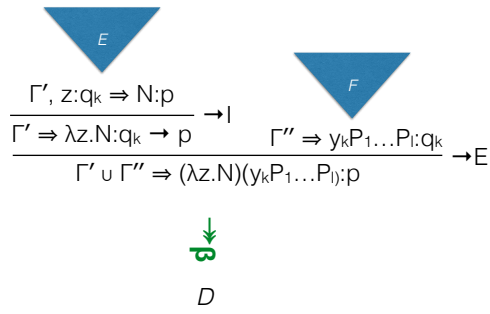
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Problematic case.

Try to extract an atomic typed subdeduction.

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Find a k such that
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 $x \in \text{FV}(N) \cap \text{FV}(y_k P_1 \dots P_l)$ implies $\Gamma(x) \in \text{At}$



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Important lemma concerning the negative subpremise property.

$$\alpha = \beta_1 \rightarrow \dots \rightarrow \beta_i \rightarrow \dots \rightarrow \beta_n \rightarrow p$$

↑
negative subpremise of α

- If $\gamma = \beta_i$ or γ is a positive subpremise of β_i , then γ is a negative subpremise of α .
- If γ is a negative subpremise of β_i , then γ is a positive subpremise of α .

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$$\Gamma = x_1:\alpha_1, \dots, x_n:\alpha_n$$

γ is a positive/negative subpremise of $\Gamma \Rightarrow \alpha$

if

γ is a positive/negative subpremise of $\alpha_1 \rightarrow \dots \rightarrow \alpha_n \rightarrow \alpha$

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Definition $\Gamma \Rightarrow \alpha$ has negative subpremise propertyfor all negative subpremises γ, δ of $\Gamma \Rightarrow \alpha$,
 $\text{tail}(\gamma) = \text{tail}(\delta)$ implies $\gamma = \delta$

negatively non-duplicated



negative subpremise property

 $((p \rightarrow q) \rightarrow r) \rightarrow s, p \rightarrow \bar{q}, (p \rightarrow q) \rightarrow r \Rightarrow s$

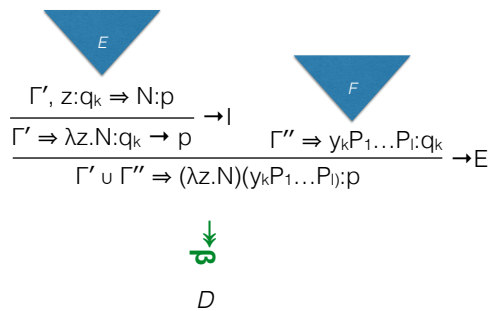
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Lemma $\Gamma \Rightarrow \alpha, \Gamma' \Rightarrow \alpha$ negatively non-duplicated $\Gamma \cup \Gamma'$ injective $\Gamma \cup \Gamma' \Rightarrow \alpha$ has negative subpremise property $\vdash \Gamma \Rightarrow M:\alpha$ $\vdash \Gamma' \Rightarrow M':\alpha$  $M =_{\beta\eta} M'$

This lemma implies Aoto and Ono's generalization of the coherence theorem.

The inductive proof is easy with "negative subpremise property".

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Find a k such that $M \equiv_{\alpha} N[z := y_k P_1 \dots P_l]$ and $x \in \text{FV}(N) \cap \text{FV}(y_k P_1 \dots P_l)$ implies $\Gamma(x) \in \text{At}$ 

The lemma is used to find this k .

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-