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Syntactic Features for Regular Constraints and an Approximation of Directional Slashes in Abstract Categorial Grammars

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Plan

- · What is ACG?
- How to obtain syntactic features expressing regular constraints
- Gapping

This talk will have a difference emphasis than my paper in the proceedings.

Lambda Caisauluscand Formal
Grammar

4th workshop

Sep. 18-19 2007, Nancy (France)

Program

An Active Control of September 18th

September 18

th Lambda Calculus and Formal Grammar Wo

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"Why is it called abstract categorial grammar?"

-Anna Chernilovskaya

Abstract Categorial Grammar (de Groote 2001, Muskens 2001)

- not a new kind of categorial grammar
- represents basic building blocks of the grammar as well as grammatical operations on them with typed linear λ -terms
- a general formalism meant to be restricted in various ways to produce more constrained grammars
- generalizes
 - CFG (context-free grammar)
 - TAG (tree-adjoining grammar)
 - MCFG (multiple context-free grammar) (or LCFRS (linear context-free rewriting system))
 - but not Lambek categorial grammar
- is like categorial grammar in that semantic composition is a homomorphic image of syntactic derivation
- treats form and meaning symmetrically

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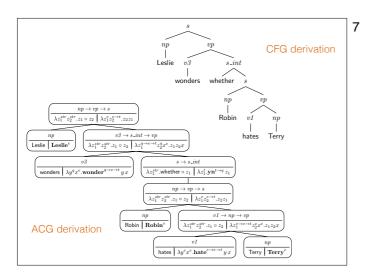
This slide is from a talk I gave in 2005. The title of my talk was "Abstract Categorial Grammar and Linear Logic".

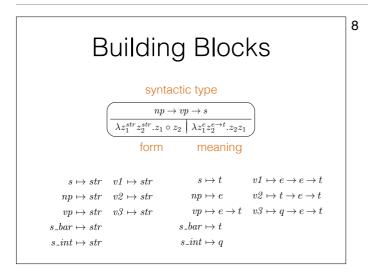
Mitsuhiro Okada introduced me saying I was going to talk about the latest exciting developments in categorial grammar.

ACG is **not** a categorial grammar.

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The connection of ACGs with context-free/tree-adjoining grammars and mainstream formal language theory is more important than the aspects of ACGs inherited from the categorial grammar tradition. "Abstract context-free grammar" or "abstract tree-adjoining grammar" would have been at least as appropriate.





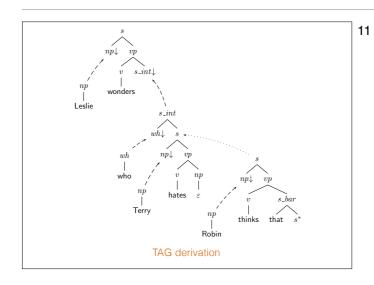
One important respect in which ACG is not a categorial grammar.

$$\frac{np \to vp \to s}{ \lambda z_1^{str} z_2^{str}. z_1 \circ z_2 \mid \lambda z_1^e z_2^{e \to t}. z_2 z_1 }$$

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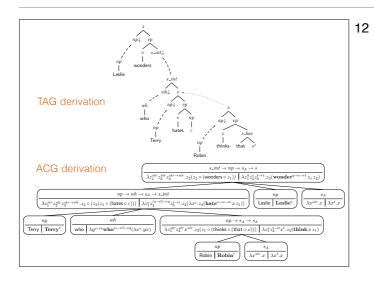
An ACG need not be lexicalized.

ACG is a generalization of TAG.



ACGs are a generalization of TAGs. Elementary trees = trees (initial trees) / unary functions on trees (auxiliary trees).

ACGs have elementary λ -terms instead of elementary trees.

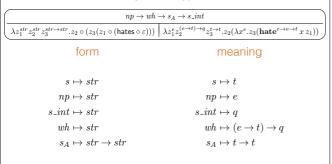


Sylvain Pogodalla has written many papers on the relation between TAGs and ACGs.

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Building Blocks

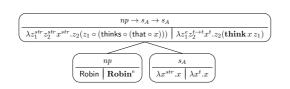
syntactic type



$$s_A \mapsto str \to str$$

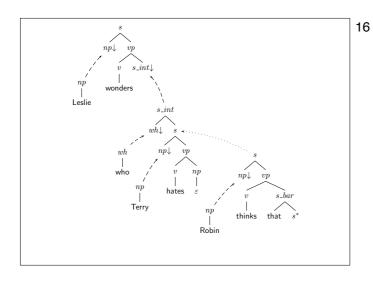
An atomic syntactic type can be mapped to a complex "prosodic" type.

Another important respect in which ACG is not a categorial grammar. The complexity of the substitution is an important parameter according to which ACGs form a hierarchy.



A constituent with a gap may have an atomic syntactic type.

An arbitrarily complex λ -term can be derived from an atomic-typed subterm of a derivation.



This illustrates a TAG-analysis of whmovement originally due to Anthony Kroch.

One could easily imagine a TAGinspired analysis of Right Node Raising in an ACG.

Gapping may be handled by a simple context-free grammar of rank ≥ 2 .

Abstract Syntactic Types

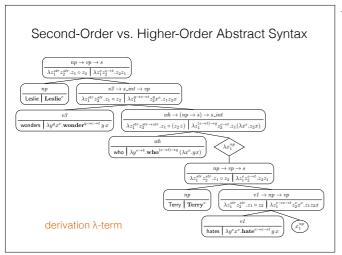
	CFG-style	TAG-style	"Lambek-style"
unsaturated standard constituents	atomic	functional	functional
modifiers	atomic	atomic	functional
continuous non-standard constituents	-	-	functional
discontinuous constituents	-	functional / atomic	functional
constituents with gaps	atomic (GPSG)	atomic	functional

You can classify various styles of analyses possible in ACGs in terms of atomic/functional distinction.

The ACG formalism supports many styles of linguistic analysis.

To take advantage of the full potential of ACGs, different styles of analyses should be explored.

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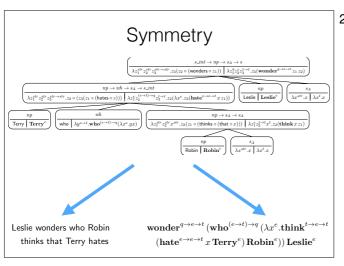
	derivations	"building block" λ-terms	language complexity	connections with computer science
second-order	trees	(almost) linear	LOGCFL	formal language theory, program schemes
		λΙ (= BCIW)	decidable	
higher-order	linear λ-terms	(almost) linear	?	BVASS
		λΙ (= BCIW)	?	

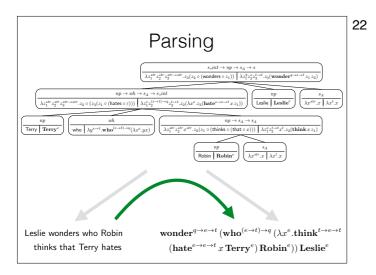
Higher-order ACGs are problematic.

Linearity is overrated.

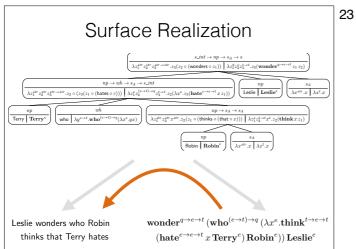
Second-order ACGs have almost no resemblance with categorial grammars.

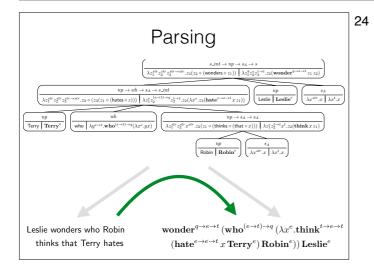
Problems that KLM pointed out have to do with higher-order ACGs. Higher-order ACGs are problematic in other ways.



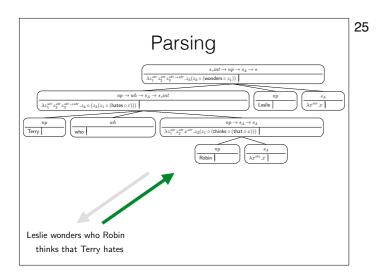


Two of the most important problems concerning grammars are parsing and surface realization (generation).

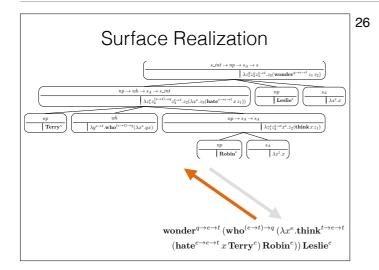




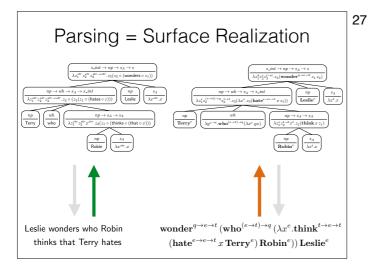
The problem boils down to finding the derivation from the input string.



The semantic component of the grammar is irrelevant.



The "form" component of the grammar is irrelevant.



In a fairly broad, interesting class of cases, the problems of parsing and surface realization have been solved.

Tabular Parsing

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?-s(0-8).
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```
\begin{split} s(x_1) &:= s_- int(x_4), np(x_3), s_A(x_1, x_2), \circ(x_2, x_3, x_4), \circ(x_4, x_5, x_6), \text{wonders}(x_5). \\ s_- int(x_1) &:= np(x_5), wh(x_2), s_A(x_3, x_4), \circ(x_1, x_2, x_3), \circ(x_4, x_5, x_6), \circ(x_6, x_7, x_8), \text{hates}(x_7), \epsilon(x_8). \\ s_A(x_1, x_8) &:= np(x_3), s_A(x_1, x_2), \circ(x_2, x_3, x_4), \circ(x_4, x_5, x_6), \circ(x_6, x_7, x_8), \text{thinks}(x_5), \text{that}(x_7). \\ s_A(x_1, x_1) &:= \\ np(x_1) &:= \text{Terry}(x_1). \\ np(x_1) &:= \text{Robin}(x_1). \\ np(x_1) &:= \text{Leslie}(x_1). \\ wh(x_1) &:= \text{who}(x_1). \end{split}
```

Tabular Realization

?-s(1).

```
\begin{split} s(x_1) &:= s\_int(x_3), np(x_4), s_A(x_1, x_2), \mathbf{wonder}(x_2, x_3, x_4). \\ s\_int(x_1) &:= np(x_5), wh(x_1, x_2, x_4), s_A(x_1, x_2), \mathbf{hate}(x_3, x_4, x_5). \\ s_A(x_1, x_3) &:= np(x_4), s_A(x_1, x_2), \mathbf{think}(x_2, x_3, x_4). \\ s_A(x_1, x_1) &:= . \\ np(x_1) &:= \mathbf{Terry}(x_1). \\ np(x_1) &:= \mathbf{Robin}(x_1). \\ np(x_1) &:= \mathbf{Leslie}(x_1). \\ wh(x_1, x_2, x_3) &:= \mathbf{who}(x_1, x_2, x_3). \end{split}
```

 $\boxed{\mathbf{wonder}(1,2,8).\ \mathbf{who}(2,3,5).\ \mathbf{think}(3,4,5).\ \mathbf{hate}(4,5,6).\ \mathbf{Terry}(6).\ \mathbf{Robin}(7).\ \mathbf{Leslie}(8).}$

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Symmetry between form and meaning is at the heart of ACG.

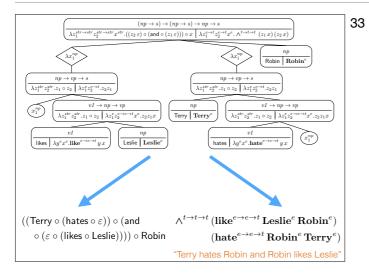


Right Node Raising

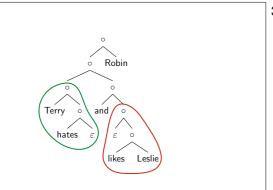
Lambek

 $((s/np)\backslash (s/np))/(s/np): \mathsf{and}: \lambda z_1^{e \to t} z_2^{e \to t} x^e. \wedge^{t \to t \to t} (z_1\, x)(z_2\, x)$

ACG



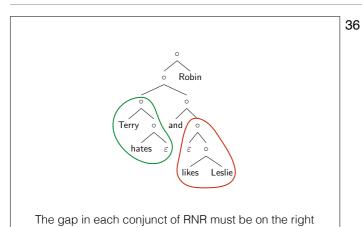
Right Node Raising.
An example of overgeneration.
The other entry considered by KLM gives "Terry hates and Robin likes Leslie".



The gap in each conjunct of RNR must be on the right periphery.

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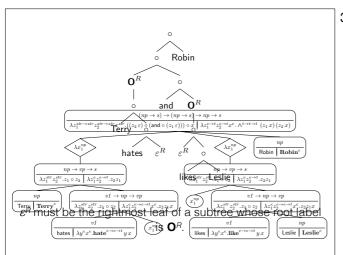
If you have a good specification, the grammar will write itself.



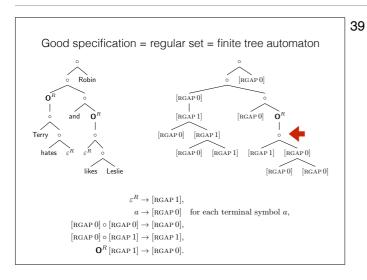
periphery.

A conjunct of RNR may contain two gaps.

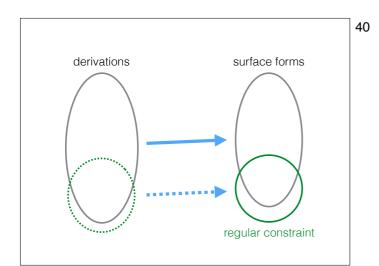
Who did [Max entice _ to read _] and [Ted ask _ to summarize _] the latest paper by Chomsky?



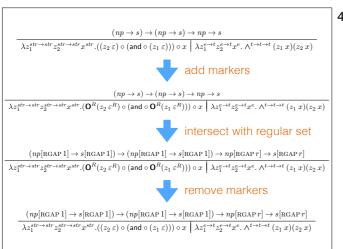
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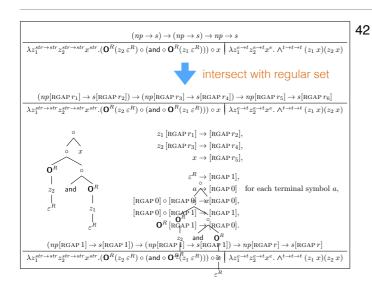


Cf. model-theoretic syntax of Jim Rogers.



A generalization of the classic result that the class of context-free languages are closed under intersection with regular sets. Building block λ -terms must be almost linear.





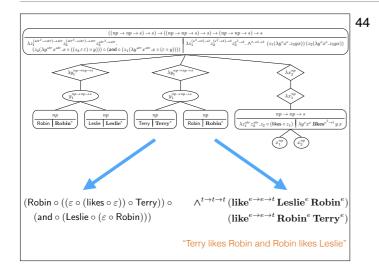
Gapping

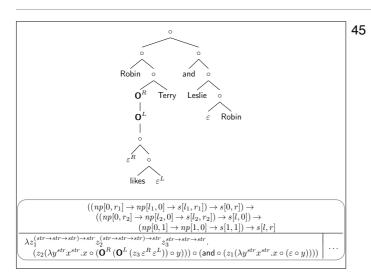
Hybrid

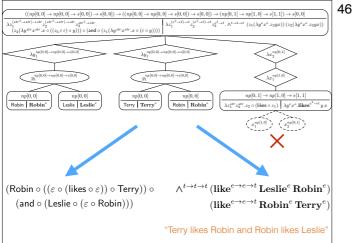
$$\frac{ \left(\left(((np \backslash s)/np) \to s \right) \to \left(((np \backslash s)/np) \to s \right) \to \left((np \backslash s)/np \right) \to s \right)}{ \lambda z_1^{str \to str} z_2^{str \to str} z_3^{str} . (z_2 z_3) \circ \left(\text{and} \circ (z_1 \varepsilon) \right) \left| \begin{array}{c} \lambda z_1^{(\varepsilon^2 \to t) \to t} z_2^{(\varepsilon^2 \to t) \to t} z_3^{\varepsilon^2 \to t}. \\ \wedge t^{\to t \to t} \left(z_1 (\lambda y^\varepsilon x^\varepsilon . z_3 y x) \right) \left(z_2 (\lambda y^\varepsilon x^\varepsilon . z_3 y x) \right) \right|$$

ACG

$$\begin{array}{c} ((np \rightarrow np \rightarrow s) \rightarrow s) \rightarrow ((np \rightarrow np \rightarrow s) \rightarrow s) \rightarrow (np \rightarrow np \rightarrow s) \rightarrow s \\ \lambda z_1^{(str^2 \rightarrow str) \rightarrow str} z_2^{(str^2 \rightarrow str) \rightarrow str} z_3^{str^2 \rightarrow str} \\ (z_2(\lambda y^{str} x^{str} . x \circ ((z_3 \varepsilon \varepsilon) \circ y))) \circ (\text{and} \circ (z_1(\lambda y^{str} x^{str} . x \circ (\varepsilon \circ y)))) \end{array}$$







The right- or left-peripherality of a gap can be enforced by a syntactic feature.

Reverse Word Order

Tom cooked the beans, and Bill, the potatoes.

Tom cooked the beans, and the potatoes, Bill.

The first remnant can be a Focus.

A. Gee, the beans and the potatoes are good! $\operatorname{Did}\nolimits$ Tom cook them again?

B1. No. Today, Tom cooked the beans, and Bill, the potatoes.

B2. No. Today, Tom cooked the beans, and the potatoes, Bill.

Topicalization + Gapping

Tom cooked the beans, and Bill, the potatoes.

Tom cooked the beans, and the potatoes, Bill.

The beans, Tom cooked, and the potatoes, Bill.

Tom cooked the beans, and the potatoes, Bill.

The last two based on Hankamer's (1979) examples.

Discontinuous Gapping

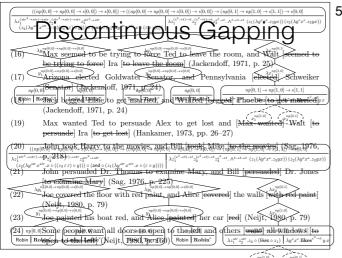
While the canonical cases of gapping have medial gaps, the gap can also be discontinuous, e.g.

Should I call you, or $_{should}$ you $_{call}$ me? Will Jimmy greet Jill first, or $_{will}$ Jill $_{greet}$ Jimmy $_{first}$? He believes her to know the answer, and she $_{believes}$ him $_{to \; know \; the \; answer}$ I expect you to help, and you $_{expect}$ me $_{to \; help}$.

-Wikepedia, Gapping

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Discontinuous Gapping

 $(((np \backslash s)/np) \to s) \to (((np \backslash s)/np) \to s) \to ((np \backslash s)/np) \to s$

- (25) Max ordered Ted to persuade Alex to get lost and [Max ordered] Walt [to persuade] Ira [to get lost]
- (26) I asked Peter to take Susan home, and [I asked] John [to take] Wendy [home]
- (27) Rarely does John call Mary at home, and [rarely does] Mary [eall] John [at home]

The type of the gap =
$$np \rightarrow np \rightarrow s$$

$$\begin{pmatrix} ((np \rightarrow np \rightarrow s) \rightarrow s) \rightarrow ((np \rightarrow np \rightarrow s) \rightarrow s) \rightarrow ((np \rightarrow np \rightarrow s) \rightarrow s) \\ \overline{\lambda z_1^{str \rightarrow str} z_2^{str \rightarrow str} z_3^{str} . (z_2 z_3) \circ (\text{and} \circ (z_1 \varepsilon))} & \overline{\lambda z_1^{(e^2 \rightarrow t) \rightarrow t} z_2^{(e^2 \rightarrow t) \rightarrow t} z_3^{e^2 \rightarrow t}}. \\ \overline{\lambda t \rightarrow t \rightarrow t} \left(z_1 (\lambda y^e x^e . z_3 y x) \right) \left(z_2 (\lambda y^e x^e . z_3 y x) \right) \end{pmatrix}$$

"I asked Peter to take Susan home and I asked Wendy to take John home."

$$\begin{array}{|c|c|c|}\hline & ((np \rightarrow np \rightarrow s) \rightarrow s) \rightarrow ((np \rightarrow np \rightarrow s) \rightarrow s) \rightarrow ((np \rightarrow np \rightarrow s) \rightarrow s) \\ \hline & \lambda z_1^{str \rightarrow str} z_2^{str \rightarrow str} z_3^{str} . (z_2 z_3) \circ (\operatorname{and} \circ (z_1 \varepsilon)) & \lambda z_1^{(e^2 \rightarrow t) \rightarrow t} z_2^{(e^2 \rightarrow t) \rightarrow t} z_2^{e^2 \rightarrow t} . \\ & \lambda z_1^{-t+t} . (z_1 (\lambda y^e x^e . z_3 y x)) (z_2 (\lambda y^e x^e . z_3 y x)) \end{array}$$

In the first conjunct of Gapping, the correspondent of the first remnant must precede the correspondent of the second remnant.

If, as has often been argued (Kuno, 1976; Neijt, 1980; Coppock, 2001; Johnson, 2014), the relative positions of the correspondents/remnants obey some (but perhaps not all) of the island constraints governing wh-extraction, those constraints can also be captured by the syntactic feature, as long as they are regular.

Conclusion

- ACG supports non-categorial-style analyses.
- In a Lambek-style analysis, any regular constraint on positions of gaps can be captured by syntactic features.